

```
public class Puzzle10 {  
  
    public static void main(String[] args) {  
        ((Puzzle10) null).print();  
    }  
  
    private static void print() {  
        System.out.println("Hello World!");  
    }  
}
```

Brainf*ck Lexical Analysis

```
MOVE_RIGHT: '>';  
MOVE_LEFT: '<';  
INCREMENT: '+';  
DECREMENT: '-';  
WRITE: '.';  
READ: ',';  
LOOP_HEADER: '[';  
LOOP_FOOTER: ']';
```

Program: **++[>+[+]].**

Program Tokens: INCREMENT INCREMENT LOOP_HEADER MOVE_RIGHT INCREMENT LOOP_HEADER INCREMENT
LOOP_FOOTER LOOP_FOOTER WRITE <EOF>

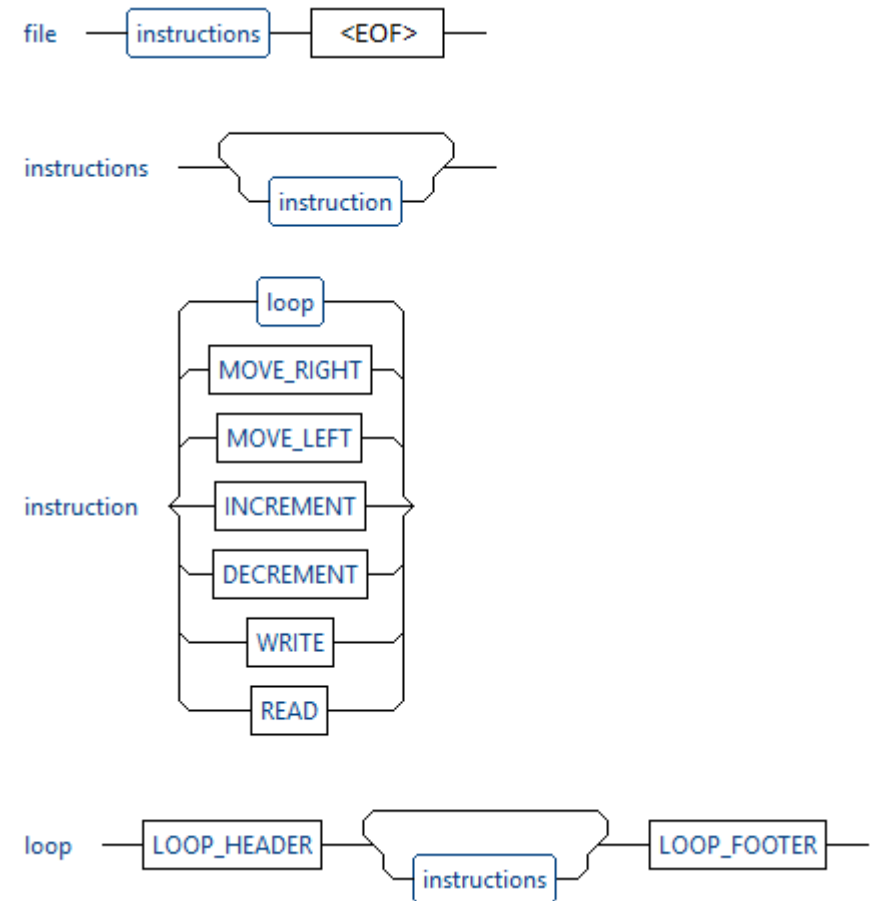
Brainf*ck Parsing Rules

```
file: instructions EOF;
```

```
instructions: instruction+;
```

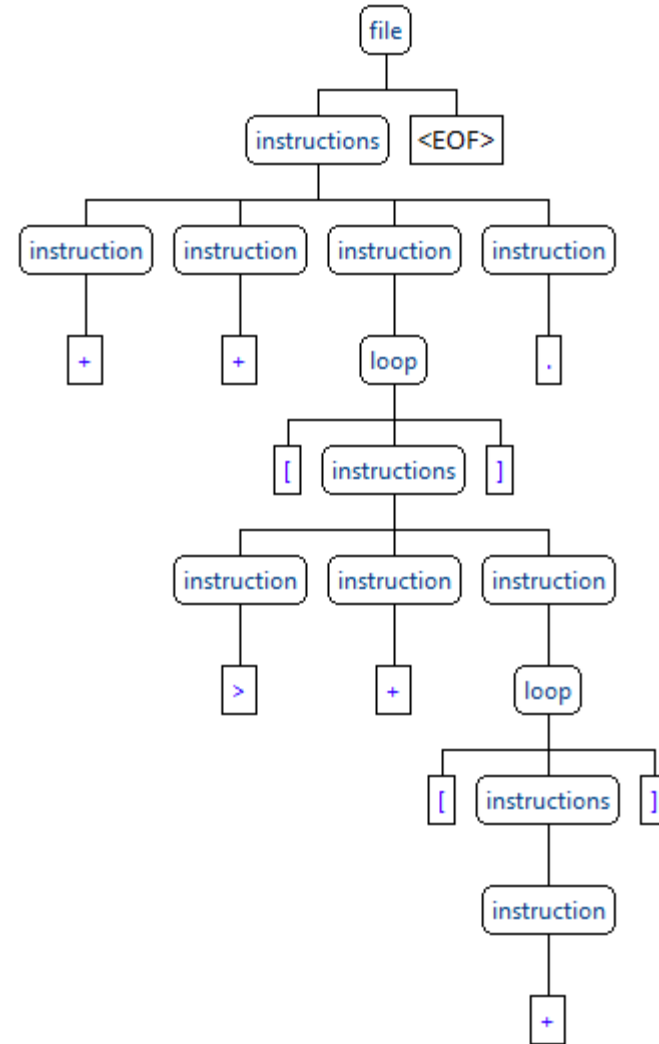
```
instruction: loop  
            | MOVE_RIGHT  
            | MOVE_LEFT  
            | INCREMENT  
            | DECREMENT  
            | WRITE  
            | READ  
            ;
```

```
loop: LOOP_HEADER instructions+ LOOP_FOOTER;
```

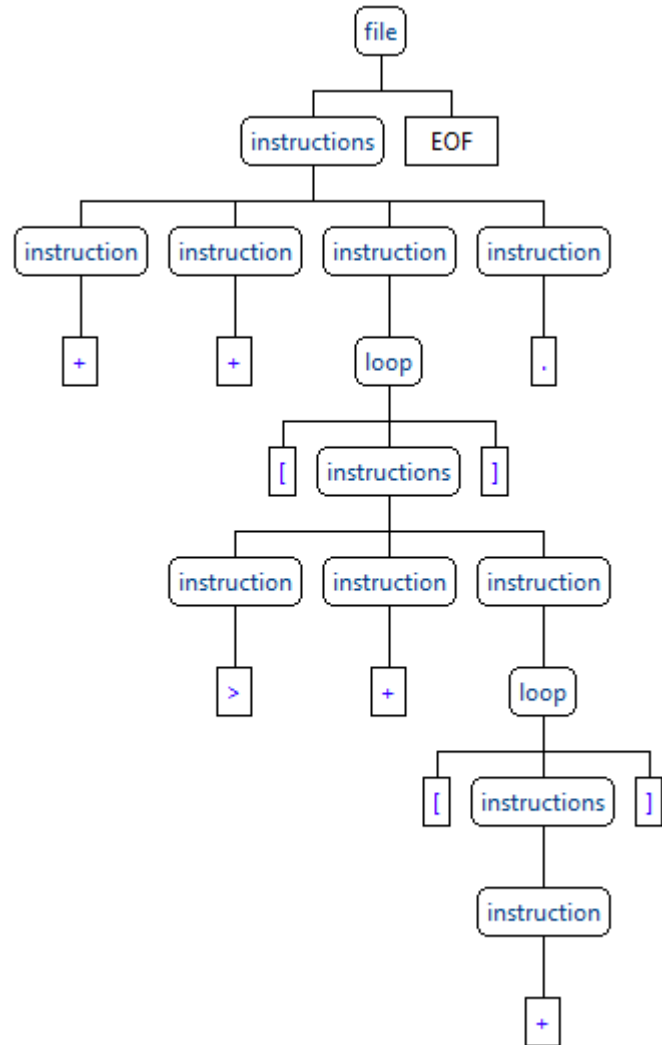


Brainf*ck Parse Tree

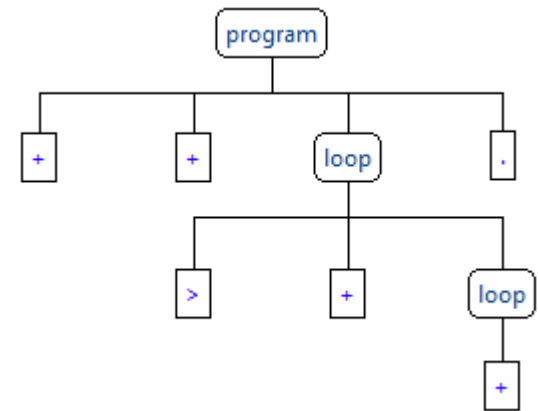
Program: **++[>+[+]].**



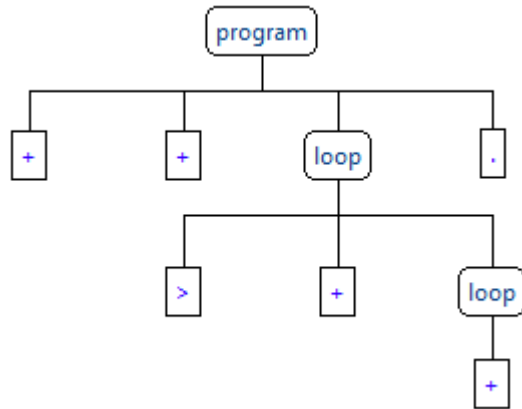
Brainf*ck Abstract Syntax Tree (AST)



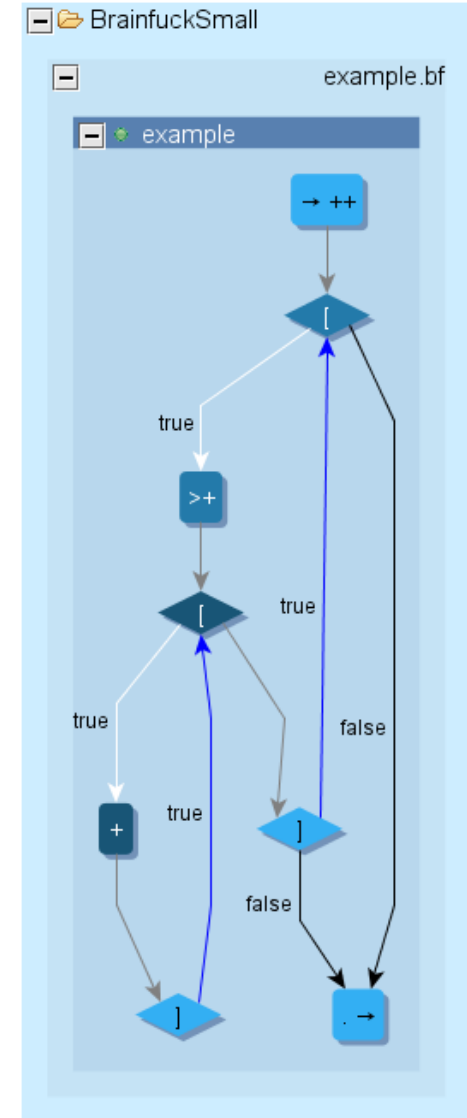
Parse Tree(s) to AST



Brainfuck*ck AST to Program Graph



Parse Tree(s) to AST



Elemental: A Brainf*ck Derivative

- github.com/benjholla/Elemental
 - Goal is to be basic, not to be tiny
 - Separates looping and branching
 - New features to explore impacts of modern language features

Instruction	Description
+	Increment the byte at the current tape cell by 1
-	Decrement the byte at the current tape cell by 1
<	Move the tape one cell to the left
>	Move the tape one cell to the right
,	(Store) Read byte value from input into current tape cell
.	(Recall) Write byte value to output from current tape cell
((Branch) If the byte value at the current cell is 0 then jump to the instruction following the matching), else execute the next instruction
[(While Loop) If the byte value at the current cell is 0 then jump to the instruction following the matching], else execute instructions until the matching] and then unconditionally return to the [
[0-9]+:	(Function) Declares a uniquely named function (named [0-9]+ within range 0-255)
{[0-9]+}	(Static Dispatch) Jump to a named function
?	(Dynamic Dispatch/Function Pointer) Jumps to a named function with the value of the current cell
"[0-9]+"	(Label) Sets a unique label (named [0-9]+ within range 0-255) within a function
'[0-9]+'	(GOTO) Jumps to a named label within the current function
&	(Computed GOTO) Jumps to the named label within the current function with the value of the current cell
#	A one line comment

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Elemental: A Brainf*ck Derivative

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 - Goal is to be basic, not to be tiny
 - Separates looping and branching
 - New features to explore impacts of modern language features
- '?' could pass control to any function!
- '&' could jump to any line!
- Goto labels with '?' or '&' could be simulated with branching or loops
- These blur control flow with data

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Positive Trend – Addressing the Languages

- Data drives execution
 - Data is half of the program!
 - “The illusion that your program is manipulating its data is powerful. But it is an illusion: The data is controlling your program.”
- Crema: A LangSec-Inspired Programming Language
 - Giving a developer a Turing complete language for every task is like giving a 16 year old a formula one car (something bad is bound to happen soon)

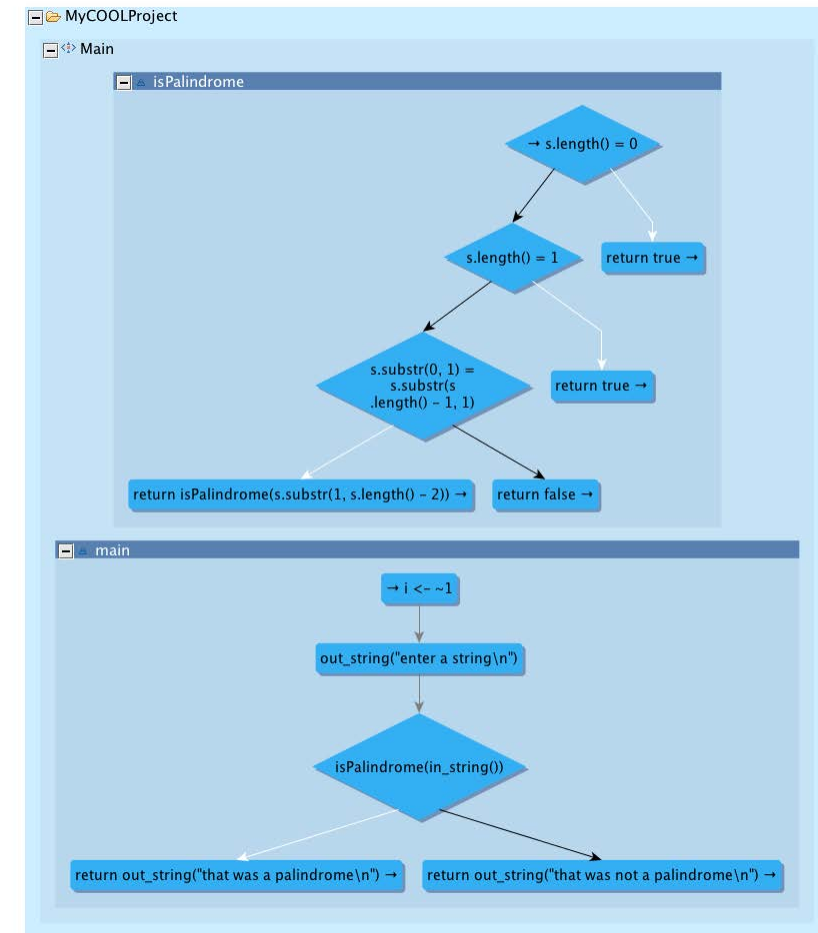


Positive Trend – Addressing the Languages

- Data drives execution
 - Data is half of the program!
 - “The illusion that your program is manipulating its data is powerful. But it is an illusion: The data is controlling your program.”
- Crema: A LangSec-Inspired Programming Language (DARPA Pilot Study)
 - Giving a developer a Turing complete language for every task is like giving a 16 year old a formula one car (something bad is bound to happen soon)
 - Apply principle of least privilege to computation (least computation principle)
 - Computational power exposed to attacker *is* privilege. Minimize it.
 - Try copy-pasting the XML billion-laughs attack from Notepad into MS Word if you want to see why...

Scaling Up: Program Analysis for COOL

- Classroom Object Oriented Language (COOL)
 - [https://en.wikipedia.org/wiki/Cool_\(programming_language\)](https://en.wikipedia.org/wiki/Cool_(programming_language))
 - <http://openclassroom.stanford.edu/MainFolder/CoursePage.php?course=Compilers>
- COOL Program Graph Indexer
 - Type hierarchy
 - Containment relationships
 - Function / Global variable signatures
 - Function Control Flow Graph
 - Data Flow Graph (in progress)
 - Inter-procedural relationships:
 - Call Graph (implemented via compliance to XCSG!)
 - <https://github.com/benjhollo/AtlasCOOL> (currently private)



Program Analysis for Contemporary Languages

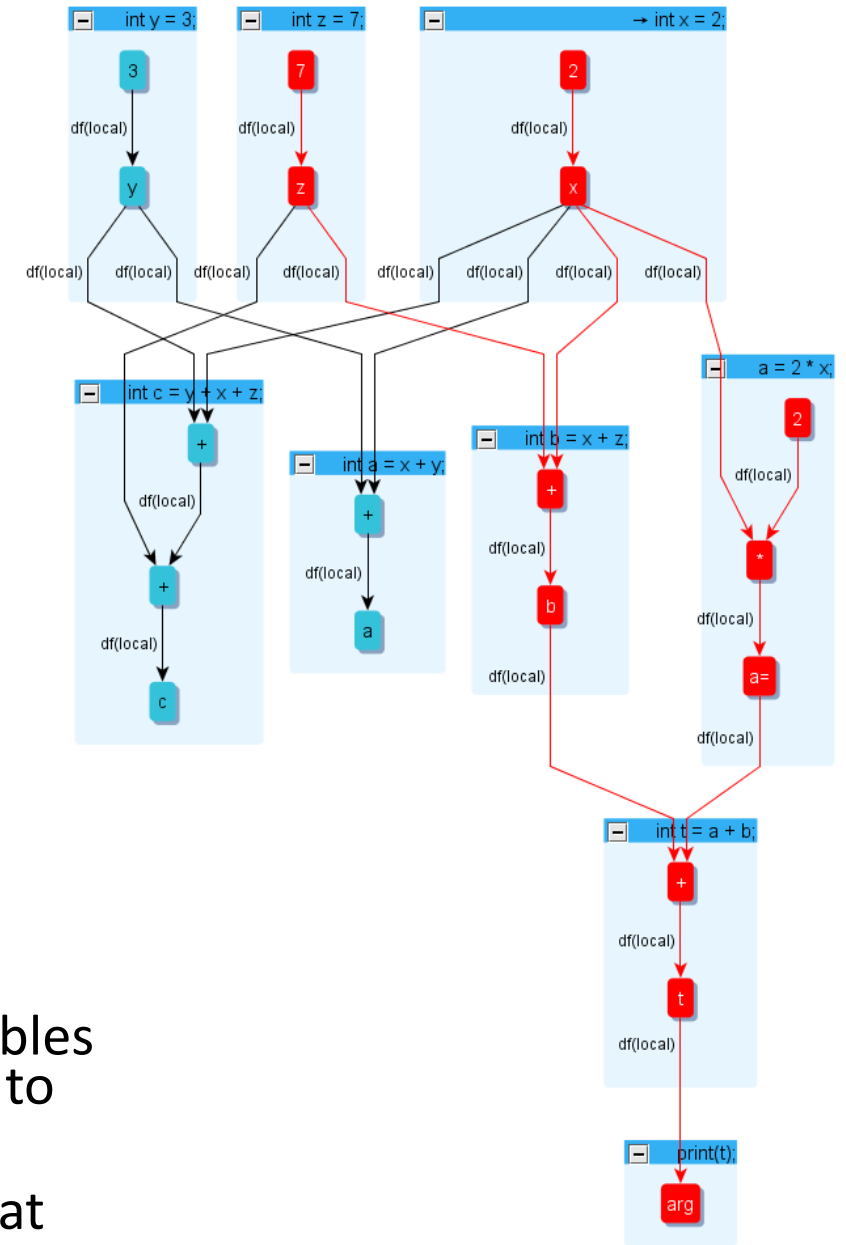
- <http://www.ensoftcorp.com/atlas> (Atlas)
 - C, C++, Java Source, Java Bytecode, *and now Brainfuck/COOL!*
- <https://scitools.com> (Understand)
 - C, C++ Source
- <http://mlsec.org/joern> (Joern)
 - C, C++, PHP Source
- <https://www.hex-rays.com/products/ida> (IDA)
- <https://binary.ninja> (Binary Ninja)
- <https://www.radare.org> (Radare)

Data Flow Graph (DFG)

Example:

1. `x = 2;`
2. `y = 3;`
3. `z = 7;`
4. `a = x + y;`
5. `b = x + z;`
6. `a = 2 * x;`
7. `c = y + x + z;`
8. `t = a + b;`
9. `print(t);` ← detected failure

Relevant lines:
1,3,5,6,8

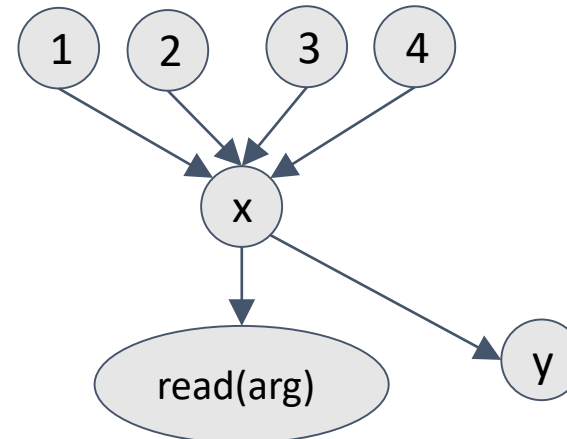


What lines must we consider if the value of `t` printed is incorrect?

- A *Data Flow Graph* creates a graph of primitives and variables where each assignment represents an edge from the RHS to the LHS of the assignment
- The *Data Flow Graph* represents global data dependence at the operator level (the atomic level) [FOW87]

Code Transformation (before – flow insensitive): Static Single Assignment Form

1. `x = 1;`
2. `x = 2;`
3. `if(condition)`
4. `x = 3;`
5. `read(x);`
6. `x = 4;`
7. `y = x;`



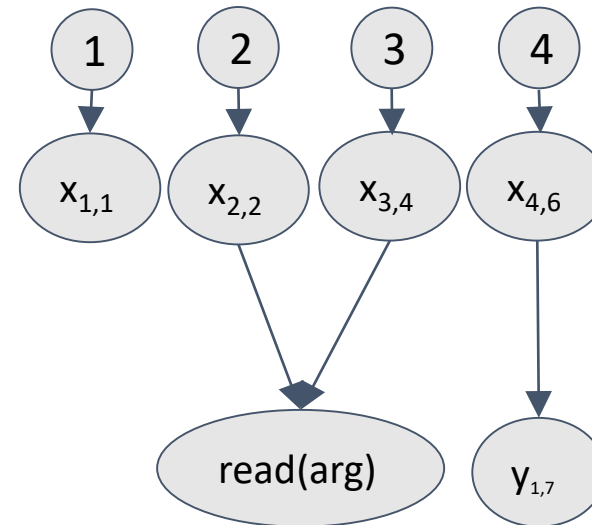
Resulting graph when statement ordering is not considered.

Code Transformation (after – flow sensitive): Static Single Assignment Form

1. $x = 1;$
2. $x = 2;$
3. $\text{if}(\text{condition})$
4. $x = 3;$
5. $\text{read}(x);$
6. $x = 4;$
7. $y = x;$



1. $x_{1,1} = 1;$
2. $x_{2,2} = 2;$
3. $\text{if}(\text{condition})$
4. $x_{3,4} = 3;$
5. $\text{read}(x_{2,2,3,4});$
6. $x_{4,6} = 4;$
7. $y_{1,7} = x_{4,6};$



Note: <Def#,Line#>

Points-to (Pointer) Analysis

- Could we answer whether or not two variables point-to the same value in memory?
- Why do we even care?
 - “Virtually all interesting questions one may want to ask of a program will eventually need to query the possible values of a pointer expression, or its relationship to other pointer expressions.”
 - Constant propagation
 - Precise call graph construction
 - Dead code elimination
 - Immutability analysis
 - Etc.

Points-to Analysis

- Could we answer whether or not two variables *may* point-to the same value in memory?
- Could we answer whether or not two variables *must* point-to the same value in memory?

Points-to Analysis

- Easy (useless) Solution:
 - A variable *must* at least point-to nothing (null)
 - Every variable *may* at most point-to anything
- Perfect (impossible) Solution:
 - A perfect Points-to is undecidable [Landi1992] [Ramalingan1994]

Andersen-style Points-to Analysis

- Flow-insensitive
 - The order of statements is not considered (does not leverage control flow graph)
- Analysis
 1. Identify each memory value to track
 2. Consider pointer assignments as subset constraints

Constraint type	Assignment	Constraint	Meaning
Base	$a = \&b$	$a \supseteq \{b\}$	$\text{loc}(b) \in \text{pts}(a)$
Simple	$a = b$	$a \supseteq b$	$\text{pts}(a) \supseteq \text{pts}(b)$
Complex	$a = *b$	$a \supseteq *b$	$\forall v \in \text{pts}(b). \text{pts}(a) \supseteq \text{pts}(v)$
Complex	$*a = b$	$*a \supseteq b$	$\forall v \in \text{pts}(a). \text{pts}(v) \supseteq \text{pts}(b)$

Andersen-style Points-to Analysis

- Fixed-point Algorithm Sketch (for Java)
 1. Identify each value to track (i.e. “new” → XCSG.Instantiation) and assign it a unique “address”
 2. Create a worklist of nodes with addresses to propagate and initialize with each addressed node
 3. If the worklist is not empty, remove a node from the worklist
 - Propagate the addresses of the node to each data flow successor node
 - If the data flow successor node received new addresses then add the successor node to the worklist
 - Repeat step 3
 4. When the algorithm reaches a fixed-point (no addresses left to propagate) then the points-to sets have been computed

Andersen-style Points-to Analysis

- Worst Case Performance?
- Worst Case: Every variable is assigned to every other variable.
 - This is the handshake problem $\rightarrow n * (n-1) \rightarrow O(n^2)$ for each iteration
 - Statements are being processed out of order, so processing a new statement could cause you to redo all previous work $\rightarrow n * (n^2) \rightarrow O(n^3)$